Principles of Operating Systems CS 446/646

5. Input/Output

- a. Overview of the O/S Role in I/O
- **b.** Principles of I/O Hardware

c. I/O Software Layers

- ✓ Overview of the I/O software
- ✓ Interrupt handlers
- ✓ Device drivers
- ✓ Device-independent I/O software
- ✓ User-level I/O system calls
- d. Disk Management

5.c I/O Software Layers Overview of the I/O software

➢ Goals and services of the I/O software

✓ device independence

- access any new I/O device without rewriting the O/S
- ✓ uniform naming
 - abstract naming space independent from physical device

✓ error handling

- lower layers try to handle the error before upper levels
- ✓ asynchronous transfers
 - make interrupt-driven operations look blocking to processes

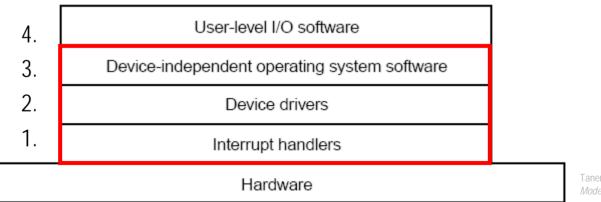
✓ buffering

decouple transfer rates and insulate data from swapping

Overview of the I/O software

> The I/O component of the O/S is organized in layers

- 1. interrupt handlers
- 2. device drivers
- 3. device-independent I/O
- 4. user-level I/O system calls



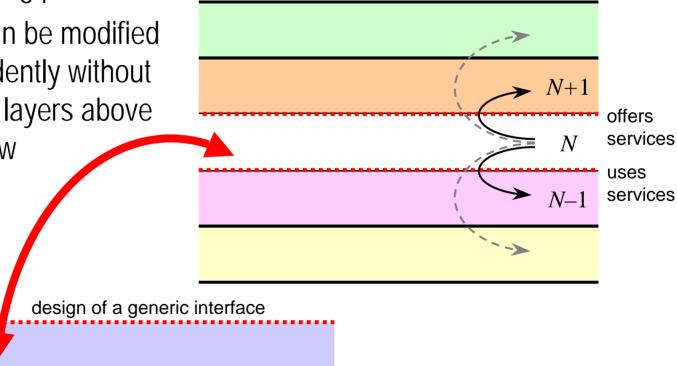
Tanenbaum, A. S. (2001) Modern Operating Systems (2nd Edition).

Typical layers of the I/O software subsystem

Overview of the I/O software

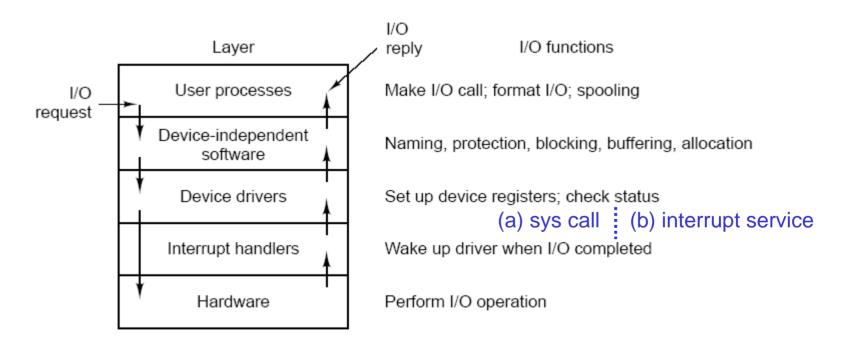
Abstraction, encapsulation and layering

- any complex software \checkmark engineering problem
- layers can be modified \checkmark independently without affecting layers above and below



Overview of the I/O software

Typical flow of control through the I/O layers upon an I/O request



Tanenbaum, A. S. (2001) *Modern Operating Systems (2nd Edition).*

1. Interrupt handler routines

- ✓ interrupts basically use the same mechanism as exceptions and traps
- ✓ when an interrupts happen, the CPU saves a small amount of state and jumps to an interrupt-handler routine at a fixed address in memory
- ✓ the interrupt routine's location is determined by an interrupt vector

1. Interrupt handler routines (cont'd)

	vector number	description
	0	divide error
	1	debug exception
	2	null interrupt
	3	breakpoint
	4	INTO-detected overflow
	5	bound range exception
	6	invalid opcode
	7	device not available
nonmaskable,	8	double fault
used for various error conditions	9	coprocessor segment overrun (reserved)
	10	invalid task state segment
	11	segment not present
	12	stack fault
	13	general protection
	14	page fault
	15	(Intel reserved, do not use)
	16	floating-point error
	17	alignment check
	18	machine check
maskable, used for $^{-1}$	19Đ31	(Intel reserved, do not use)
device-generated	32Ð255	maskable interrupts
interrupts	Intel Pentium processor event-vector table Silberschatz, A., Galvin, P. B. and Gagne Operating Systems Concepts with Java (

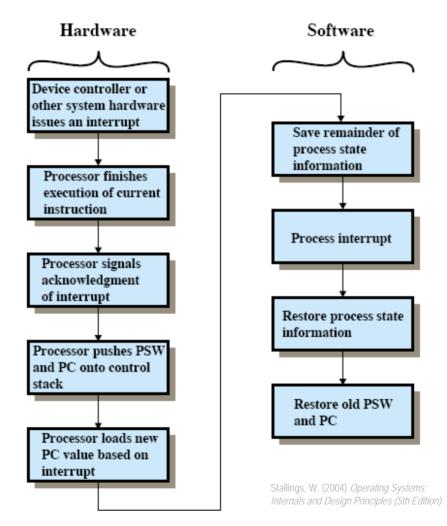
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CS 446/646 - Principles of Operating Systems - 5. Input/Output

Operating Systems Concepts with Java (6th Edition,

1. Interrupt handler routines

- ✓ typical steps followed by an interrupt routine:
 - a. save any registers not saved by the interrupt hardware
 - b. set up a context (TLB, MMU, page table) for the routine
 - c. set up a stack for the routine
 - d. acknowledge the interrupt controller
 - e. extract information from the I/O device controller's registers
 - f. etc.
- ✓ interrupt processing is a complex operation that takes a great number of CPU cycles, especially with virtual memory



Simple interrupt processing

2. Device drivers

- ✓ each I/O device needs a device-specific code to control it
- ✓ device manufacturers supply drivers for several popular O/S
- ✓ a driver handles one type of device or one class (ex: SCSI)
- ✓ the driver logic is generally executed in kernel space (although microkernel architectures might push it in user space)
- ✓ drivers should "snap into place" in the kernel through deviceindependent interfaces (see next section)
- \checkmark two main categories of drivers
 - block-device drivers: disks, etc.
 - character-device drivers: keyboards, printers, etc.

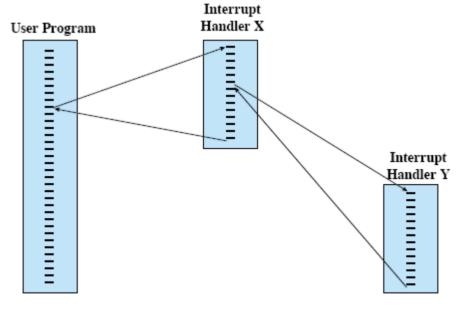
2. Device drivers (cont'd)

- \checkmark a driver has several functions
 - accept abstract read/write requests from the deviceindependent software above and translate them into concrete I/O-module-specific commands
 - initialize the device, if needed
 - manage power requirements
 - Iog device events

2. Device drivers (cont'd)

- \checkmark typical code organization of a device driver:
 - a. check validity of input parameters coming from above
 - b. if valid, translate to concrete commands, e.g., convert block number to head, track & sector in a disk's geometry
 - c. check if device currently in use; if yes, queue request; if not, possibly switch device on, warm up, initialize, etc.
 - d. issue appropriate sequence of commands to controller
 - e. if needs to wait, block
 - f. when interrupted, check for errors and pass data back
 - g. process next queued request

- 2. Device drivers (cont'd)
 - ✓ a driver code must be reentrant to allow for nested interrupts



(b) Nested interrupt processing

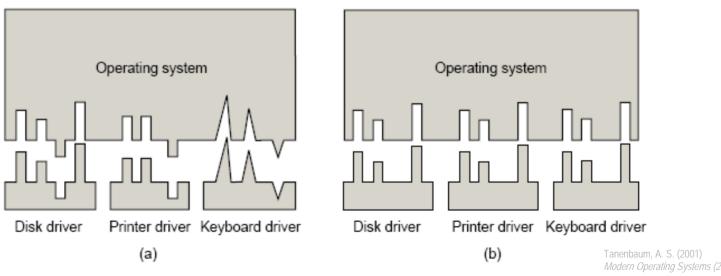
Stallings, W. (2004) Operating Systems:

5.c I/O Software Layers Device-independent I/O software

- 3. Device-independent I/O software
 - ✓ generic functions provided by the kernel I/O subsystem:
 - uniform interfacing for device drivers
 - buffering
 - error reporting
 - providing a device-independent block size

Device-independent I/O software

- 3. Device-independent I/O software (cont'd)
 - ✓ uniform interfacing
 - make all I/O devices look more or less the same, so that the O/S doesn't need to be hacked every time a new device comes along



(a) Without and (b) with a standard driver interface

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Device-independent I/O software

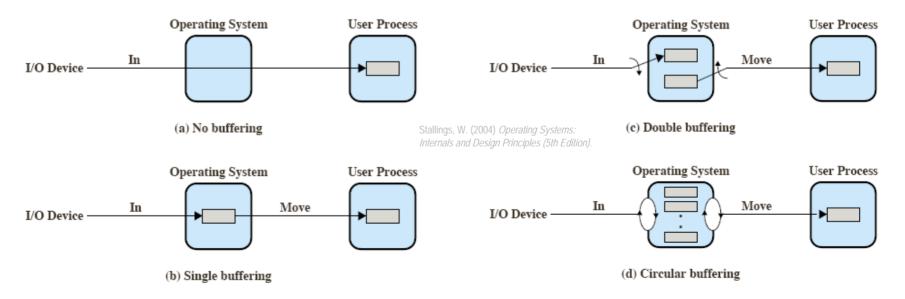
- 3. Device-independent I/O software (cont'd)
 - ✓ uniform interfacing
 - therefore, generally one unified interface
 - possibly additional specialized extensions for the main device categories
 - block devices: read(), write()
 - random-access block devices: seek()
 - character-stream devices: get(), put()
 - network devices: network socket interface

Device-independent I/O software

- 3. Device-independent I/O software (cont'd)
 - ✓ buffering = "decoupling"
 - memory area that stores data in kernel space while transferred between device and application
 - cope with a speed mismatch between producer and consumer (ex: modem thousand times slower than disk)
 - adapt between services with different data-transfer sizes (ex: fragmentation and reassembly of network packets)
 - "copy semantics": cache data while transferred so it is not affected by changes from application or swapping
 - read ahead (locality principle)

Device-independent I/O software

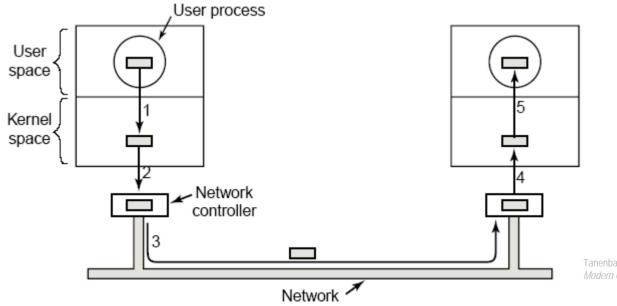
- 3. Device-independent I/O software (cont'd)
 - ✓ buffering
 - double buffering: further decouples producer from consumer (ex: modem fills 2nd buffer while 1st buffer is written to disk)
 - circular buffering: extension suitable for rapid bursts of I/O



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Device-independent I/O software

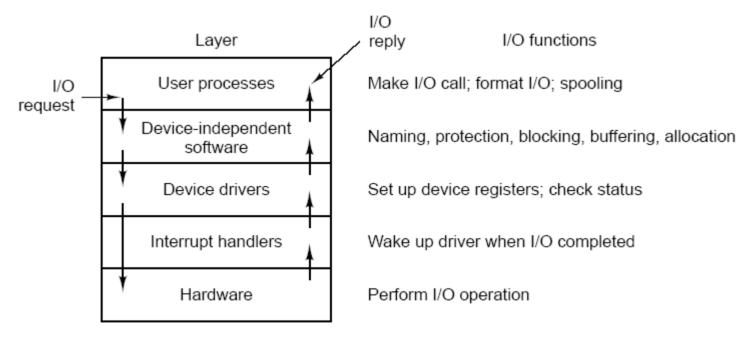
- 3. Device-independent I/O software (cont'd)
 - ✓ buffering in networking



Tanenbaum, A. S. (2001) Modern Operating Systems (2nd Edition)

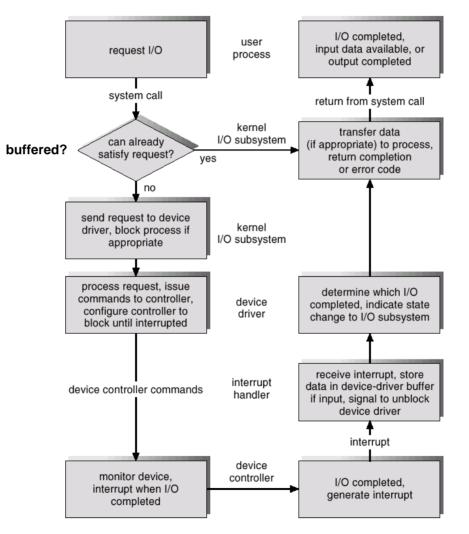
5.c I/O Software Layers User-level I/O system calls

4. User-level I/O system calls



Tanenbaum, A. S. (2001) Modern Operating Systems (2nd Edition).

5.c I/O Software Layers User-level I/O system calls



Silberschatz, A., Galvin, P. B. and Gagne. G. (2003) Operating Systems Concepts with Java (6th Edition)

The life-cycle of an I/O request

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